Non-cooperative Wireless Networks

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- With receive antenna cooperation: $C = M \log(SNR)$
- Without receive antenna cooperation: $C = \log(SNR)$, (Interference channel, Carleial, 1975)

Can we realize a pre-log of M without receive antenna cooperation?

An Interference Relay Network



Review of Previous Work

- Ground-breaking work by Gupta & Kumar, 2000 shows that capacity in large **ad-hoc networks** scales at least as $\Theta(\sqrt{n})$.
- Capacity in large **relay networks** scales as $\Theta(\log n)$ (Gastpar & Vetterli, 2002).
- Capacity in large **ad-hoc networks** with **node mobility** scales as $\Theta(n)$ (Grossglauser & Tse, 2002).
- Capacity in large **ad-hoc networks** with **network coding** scales as $\Theta(n)$ (Gupta & Kumar, 2003).
- Power efficiency in large fading relay networks scales at least as $\Theta(\sqrt{n})$ (Dana & Hassibi, 2003).

Assumptions

- All terminals are equipped with single-antenna transceivers.
- A number of designated source-destination (S-D) pairs wants to establish communication assisted by a set of relays.
- **No cooperation** between source terminals and between destination terminals.
- No direct links between source and destination terminals.
- As the **network grows large** the **number of S-D pairs** remains **constant**, the **number of relay terminals goes to infinity**.
- Encompasses traditional relay networks as special case.

MIMO Gains in Coherent Point-to-Point Links

• In an $M \times M$ coherent MIMO system, **capacity** satisfies (receive antenna cooperation necessary)

 $C \approx M \log(SNR)$

with the **multiplexing gain** given by the pre-log M.

• Array gain is the SNR improvement resulting from coherent combining. In a $1 \times M$ system with perfect receive CSI

 $C \approx \log(M \, \mathrm{SNR})$

with array gain M.

Interference Relay Network



- Source nodes transmit independent data streams.
- **Two-hop communication** using a *"listen and transmit"* protocol with *K* relays and **perfect synchronization**.
- Channels H_k and G_k are ergodic i.i.d. Gaussian block-fading.
- Random variables E_k and P_k capture **large-scale fading** and **path loss**.
- Nodes are placed in a domain of fixed area with a dead-zone around source and destination nodes ⇒ E_k and P_k are positive and bounded.

- **Destination terminals** are assumed to be able to **cooperate** and have **perfect knowledge** of the **composite MIMO channel**.
- Relay terminals have perfect knowledge of all their backward and forward channels.
- Upper bound through "max-flow min-cut theorem"

$$\sum_{i \in \mathcal{S}, j \in \mathcal{S}^c} R^{(i,j)} \leq I(\mathbf{X}^{(S)}; \mathbf{Y}^{(S^c)} | \mathbf{X}^{(S^c)}).$$

The Broadcast Cut



Cut set bound achieved if all the relay and destination terminals cooperate.

Lower Bound: Relay Partitioning



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Lower Bound: Matched Filtering



Lower bound through relay partitioning, MF, and independent decoding (i.e. no cooperation) at destination terminals.

Smart Scattering



Distributed multi-stream separation through **smart scatterers** performing **matched-filtering**.

• For $K \to \infty$, lower bound approaches upper bound and the **network** capacity converges (w.p.1) to

$$C = \frac{M}{2}\log(K) + O(1).$$

- Asymptotically in K cooperation between destination terminals is not needed to achieve network capacity.
- Independent decoding at the destination terminals achieves network capacity ⇒ significant reduction in computational complexity compared to vector decoding.

- Multiplexing gain of M/2 without cooperation between destination terminals.
- Loss in spectral efficiency (factor 1/2) due to "*listen and transmit*" protocol.
- Distributed array gain of K.

Distributed Interference Cancelation



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- Multi-stream separation realized in a completely decentralized fashion ⇒ Distributed interference cancelation.
- Network coding not needed to achieve capacity in large interference relay networks ⇒ Matched filtering is good enough.

- We relax the assumption of channel knowledge at the relays. k-th relay terminal needs to know E_k + noise variance.
- **Relays** simply perform **amplify-and-forward (AF)**.
- Receiver knows composite MIMO channel.

AF Interference Relay Networks



- In the large relay limit $K \to \infty$, **AF interference relay network approaches point-to-point MIMO system** with perfect receive CSI.
- Asymptotic capacity is half the capacity of a point-to-point coherent MIMO channel given by (receive terminal cooperation necessary)

$$C_{AF}^{\infty} = \frac{M}{2} \log(\mathrm{SNR}) + O(1).$$

• SNR depends critically on E_k .

- Multiplexing gain of M/2 realized.
- Number of relay terminals does not enter scaling law! ⇒ No distributed array gain.
- Relays can help to restore the rank of poor-scattering channels (active (but dumb) scatterers).
- Cooperation between destination terminals is crucial.

Convergence of Capacity in the AF Case



Capacity vs. number of relays for the AF interference relay network

Conclusion

- We showed that **MIMO gains** can be realized in **large interference** relay networks in a completely distributed fashion.
- Smart scatterers realize multi-stream separation without cooperation between any of the terminals.
- **Dumb scatterers** rebuild multiplexing gain in poor-scattering environments.
- **Open Issues:** Synchronization, scaling number of source-destination terminals as well.